

CS 6901 Capstone Exam Systems Fall 2017: Choose any 2 of the 3 problems.

1) Rewrite  $F(a, b, c, d) = \sum (0, 3, 4, 8, 10, 11)$  in fully simplified product-of-sums form.

2) Consider the following page reference string for a virtual memory system in which physical memory has exactly 3 frames:

2, 4, 6 4, 2, 7, 4, 3, 7, 2, 3

For each of the following page replacement algorithms, show which references will cause page faults and show the contents of the 3 frames at the time of each fault. Assume that the frames are initially empty. You do not need to show the first 3 faults that are caused by demand paging.

- a) Least Recently Used
- b) Second Chance

3) There are 3 standard goals to the 2-process mutual exclusion problem.

Goal 1: Mutual exclusion is guaranteed

Goal 2: Deadlock cannot occur.

Goal 3: Indefinite postponement cannot occur.

Attempted Solution: common variable: lock (initially false)

Assume the existence of an atomic (non-interruptible) test\_and\_set function that both returns the value of its boolean argument and sets the argument to true.

```
Process 1                               Process 2
while (true) {                           while (true) {
  while (test_and_set(lock));             while (test_and_set(lock));
  Critical section;                       Critical section;
  lock = false;                           lock = false;
  Noncritical section;                    Noncritical section;
}
```

For the above solution,

a) Select one goal that is not satisfied and provide an execution sequence that violates the goal.

b) Select one goal that is satisfied and give a brief explanation that justifies why the goal is met for all possible execution sequences.

CS 6901 Capstone Exam Data Structures and Algorithms Fall 2017

Choose any 2 of the 3 problems.

1) Given a possibly empty binary tree, write a function that returns the number of nodes in the tree that have a right child, but no left child. The prototype for your function is

```
int RightNoLeft(TreeNode *ptr) .
```

Global variables may not be used. No additional functions may be defined. Declare all data structures.

2) Given an array of  $n$  nonzero real numbers  $a[0] \dots a[n-1]$ , write a function to partition the array (not sort) so that all its negative elements come before all its positive elements. Your algorithm should have  $O(n)$  time complexity. The function prototype is

```
void negpospartition(float a[], int n) .
```

3) Count the precise number of "fundamental/basic operations" executed in the following code. Your answer should be a function of  $n$  ( $n \geq 1$ ) in closed form. Note that "closed form" means that you must resolve all  $\sum$ 's and  $\dots$ 's. An asymptotic answer (such as one that uses big-oh, big-theta, etc.) is not acceptable.

```
for(int k = 1; k < n; k++) {  
    Perform 1 fundamental/basic operation;  
    for (int j = k; j <= n; j++)  
        Perform 1 fundamental/basic operation;  
    //endfor j  
} //endfor k
```

# Theory Exam

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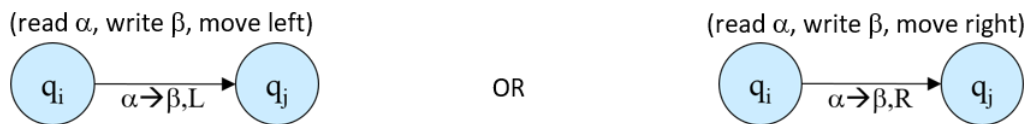
1. Give regular expressions describing each of the following languages over  $\Sigma = \{0, 1\}$ :
  - a.  $\{w : w \text{ contains the substring } 101\}$
  - b.  $\{w : w \text{ contains at least three } 0\text{'s}\}$
  - c.  $\{w : w \text{ contains at most three } 1\text{'s}\}$
  - d.  $\{w : |w| \geq 3\}$
  - e.  $\{w : |w| \leq 3\}$
  - f.  $\{w : w \text{ starts and ends with different symbols}\}$
  - g.  $\{w : \text{every odd position of } w \text{ is } 0\}$
  - h.  $\{w : w \text{ does not contain exactly two } 1\text{'s}\}$
  - i.  $\{w : \text{every } 0 \text{ in } w \text{ is followed by two } 1\text{'s}\}$
  - j.  $\{w : w \text{ starts with the substring } 011 \text{ and contains the substring } 110\}$

Each question is worth 2 points. No partial credit will be given.

2. Give the state diagram for a Turing machine that decides the following language over  $\Sigma = \{0, 1\}$ :

$$L = \{w : w \text{ contains an even number of occurrences of the substring } 011 \text{ and } |w| \text{ is odd}\}$$

Use the following notation to label each of your machine's transitions:



3. A *coloring* of a graph is an assignment of colors to its nodes so that no two adjacent vertices have the same colors.

Let **COLOR** =  $\{G, k : G = (V, E) \text{ is a graph that can be colored by } k \text{ colors}\}$ .

Provide a polynomial verifier to prove that **COLOR**  $\in$  NP.